

Unit 2

DB2 and SQL

Outline of Unit 2

2.1 Overview

2.2 Data Definition

2.3 Data Manipulation

2.4 The System Catalog

2.5 Embedded SQL

2.1 Overview

Background

- Relational Model: proposed by Codd, 1970

Ref: CACM Vol. 13, No.6, "A relational model of data for large shared data banks"

	System R	INGRES
Developer	IBM San Jose Res. Lab 1974 - 1979	UC Berkeley late 1970 - early 1980
Machine	IBM System 370	DEC PDP
O. S.	VM / CMS	UNIX
Query Language	SQL	QUEL
Language Embedded	COBOL or PL/1	COBOL, PASCAL, C FORTRAN, BASIC
Commercial Product	DB2, SQL / DS	Commercial INGRES
Distributed OB	R*	Distributed INGRES
OO Extension	Starburst	POSTGRES

Relational Databases

- Definition: A *Relational Database* is a database that is perceived by its users as a collection of tables (and nothing but tables).

<e.g.> Supplier-and-Parts Database

S

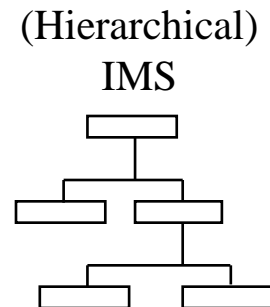
S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens

SP

S#	P#	QTY
S1	P1	300
S1	P2	200
S1	P3	400
S1	P4	200
S1	P5	100
S1	P6	100
S2	P1	300
S2	P2	400
S3	P2	200
S4	P2	200
S4	P4	300
S4	P5	400

P

P#	PNAME	COLOR	WEIGHT	CITY
P1	Nut	Red	12	London
P2	Bolt	Green	17	Paris
P3	Screw	Blue	17	Rome
P4	Screw	Red	14	London
P5	Cam	Blue	12	Paris
P6	Cog	Red	19	London



Relational Databases (cont.)

S

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens

P

P#	PNAME	COLOR	WEIGHT	CITY
P1	Nut	Red	12	London
P2	Bolt	Green	17	Paris
P3	Screw	Blue	17	Rome
P4	Screw	Red	14	London
P5	Cam	Blue	12	Paris
P6	Cog	Red	19	London

SP

S#	P#	QTY
S1	P1	300
S1	P2	200
S1	P3	400
S1	P4	200
S1	P5	100
S1	P6	100
S2	P1	300
S2	P2	400
S3	P2	200
S4	P2	200
S4	P4	300
S4	P5	400

- S, P, SP: 3 relations (tables)
- A row in a relation is called a tuple (record)
- S, P: entities; SP: relationship
- primary key: S# in S, P# in P, (S#, P#) in SP
- atomic: not a set of values, instead of repeating group

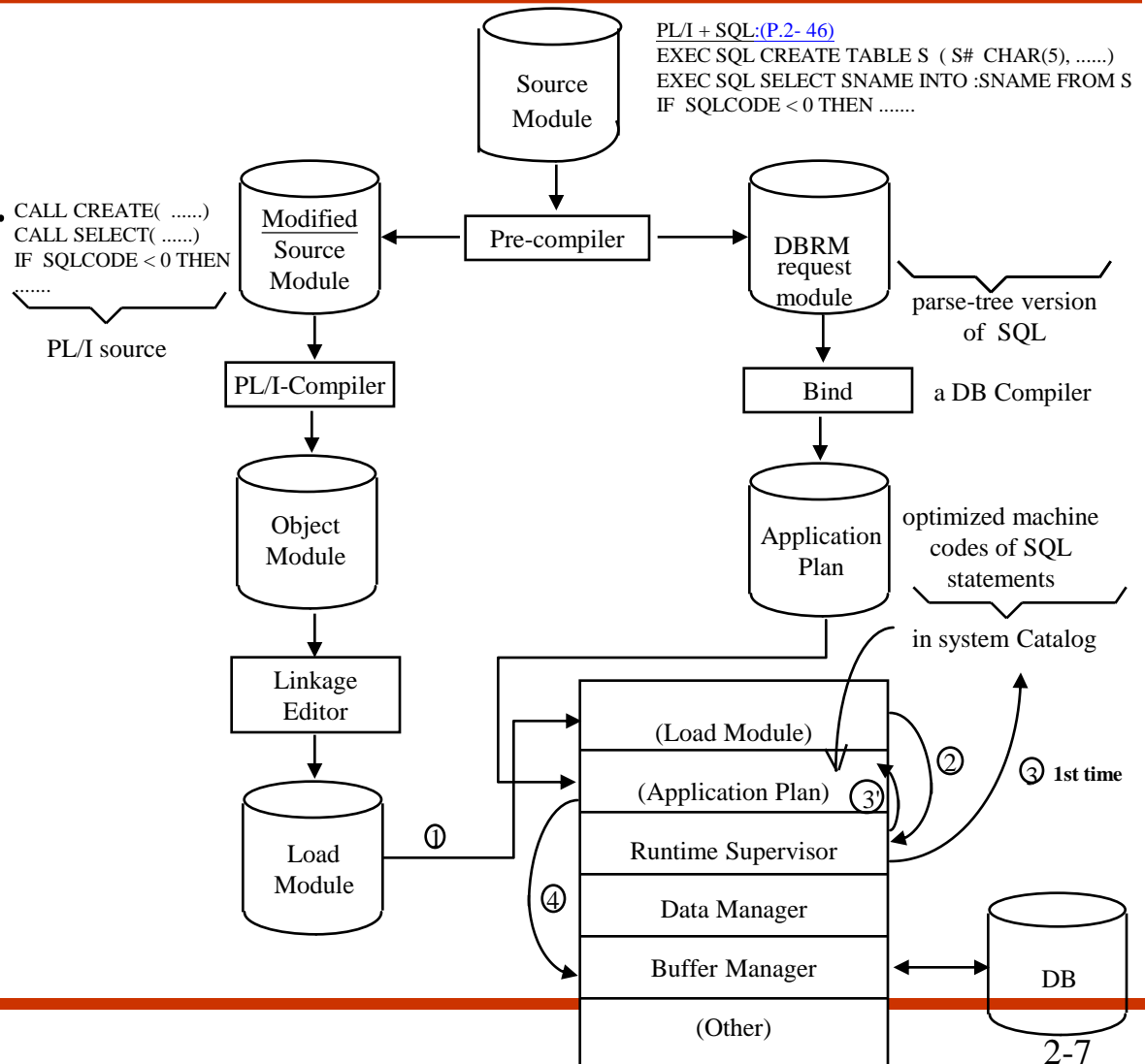
< e.g. >

S#	P#
S1	{ P1, P2, P3, P4, P5, P6 }
S2	{ P1, P2 }
⋮	⋮

⇒ atomic
Normalization

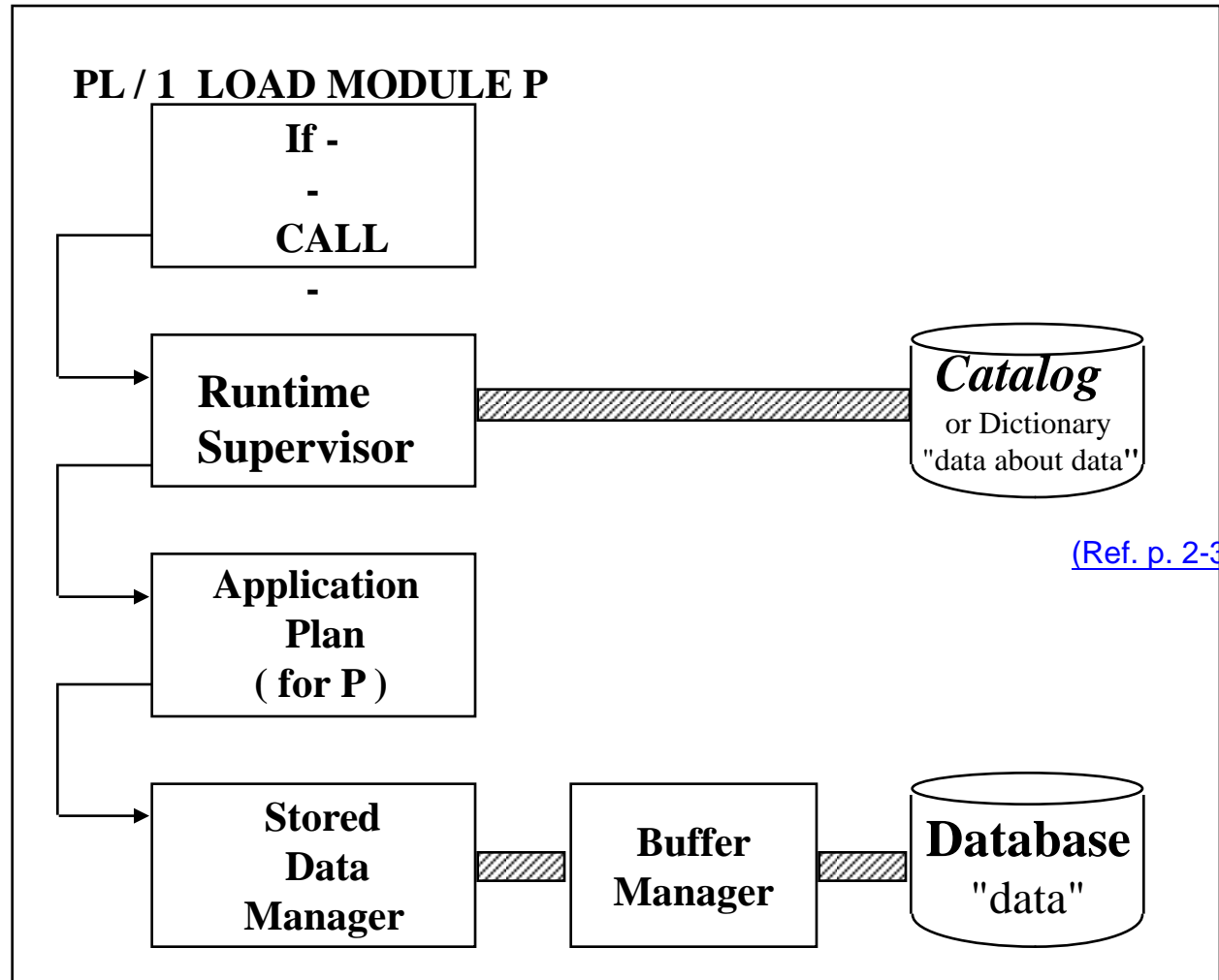
Major System Components: DB2

1. *Pre-compiled*
2. *Bind*
3. *Runtime Supervisor*
4. *Data Manager*
5. *Buffer Manager*

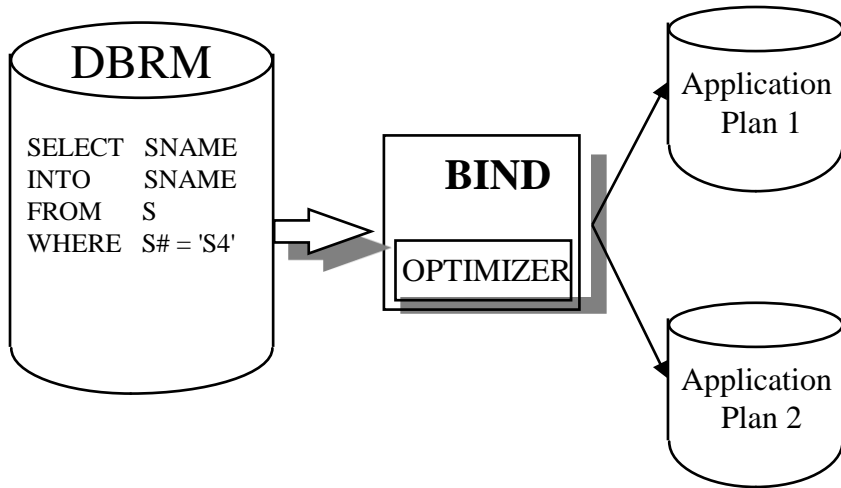


Major System Components: Execution time

Execution time



Major System Components: Optimizer



- Plan 1 (without index): **SCAN S**
 if S # = 'S4' then extract name field
 go to SCAN S

⋮

- Plan 2 (with sorted index): **Binary Search X**
 if X.key = 'S4' then

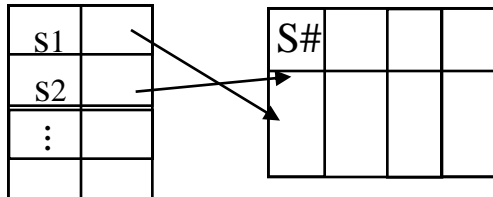
Be chosen by optimizer ([Ref. p. 2-34](#))

- Considerations :

1. Which table ?
2. How big ?
3. What index exist ?

⋮

X: S#_index



2.2 Data Definition

DDL of SQL

- **Base Table**
 - **Create Table**
 - **Drop Table**
 - **Alter Table**

- **Index**
 - **Create Index**
 - **Clustering Index**
 - **Drop Index**

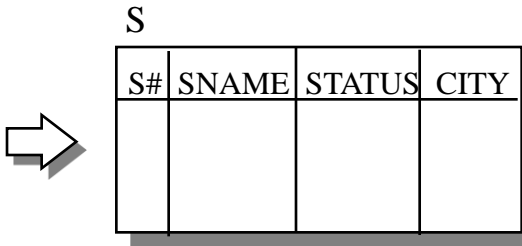
- **View**
 - **Create View**
 - **Drop View**

Base Tables

□ A named table

<e.g. 1> **CREATE TABLE S**
(S# CHAR (5) NOT NULL,
SNAME CHAR (20),
STATUS SMALLINT,
CITY CHAR (15));

S



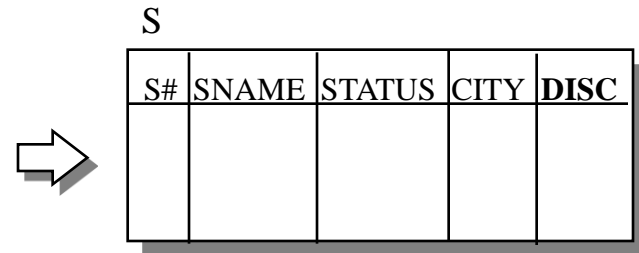
S#	SNAME	STATUS	CITY

• Data can be entered by

- (1) **INSERT** statement
- (2) DB2 load utility

<e.g. 2> **ALTER TABLE S**
ADD **DISC** SMALLINT;

S



S#	SNAME	STATUS	CITY	DISC

<e.g. 3> **DROP S ;**

- Description of S is removed from **system catalog**.
- All index and views defined on S are automatically dropped.

Base Tables (cont.)

- FOREIGN KEY

FOREIGN KEY (*column-commalist*)

REFERENCES *base-table* [(*column-commalist*)]

[ON DELETE *option*]

[ON UPDATE *option*]

- CHECK

CHECK (*conditional-expression*)

S

S#			
S ₁	Smith	

<e.g.> : CREATE TABLE SP

(Ref. p. 3-9)

(S# S# NOT NULL, P# P# NOT NULL, QTY QTY NOT NULL,

PRIMARY KEY (S#, P#),

FOREIGN KEY (S#) REFERENCES S

ON DELETE CASCADE

ON UPDATE CASCADE ,

FOREIGN KEY (P#) REFERENCES P

ON DELETE CASCADE

ON UPDATE CASCADE ,

CHECK (QTY >0 AND QTY <5001)); (Ref. p. 10-16, Integrity Rule)

SP

S#	P#	QTY
S ₁	P ₁

Indexes

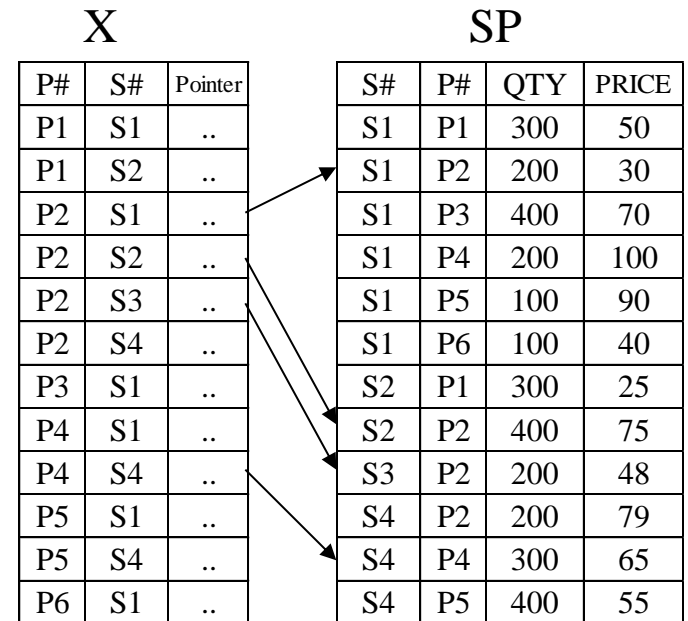
<e.g.1>: **CREATE INDEX X ON SP (P# ASC, S# ASC);**

<e.g.2> : **CREATE UNIQUE INDEX XS ON SP (S#, P#)**

- enforced that no two tuples have the same index field.

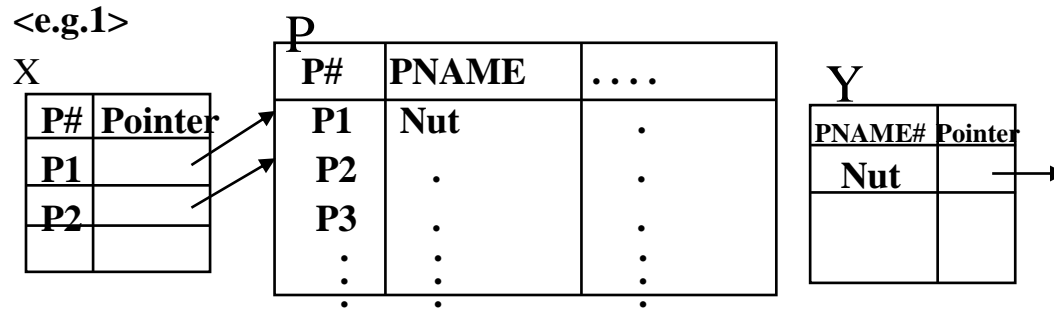
<e.g.3>: **DROP INDEX X;**

- Definition of **X** is removed from **Catalog**.



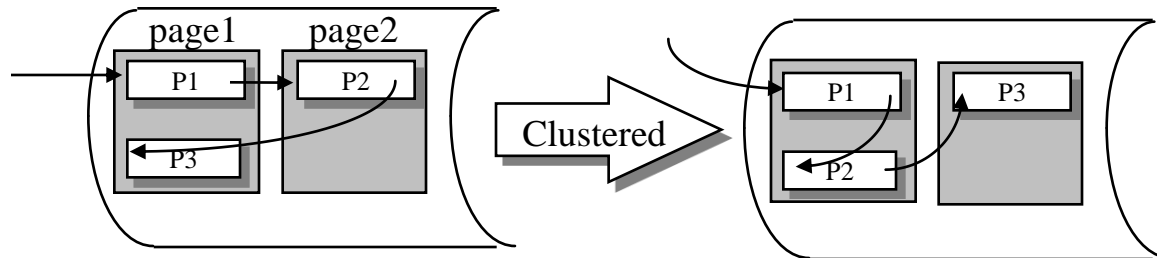
Clustering index

- Logical sequence \cong Physical sequence



logical sequence : P1 < P2 < P3 < ...

physical sequence :

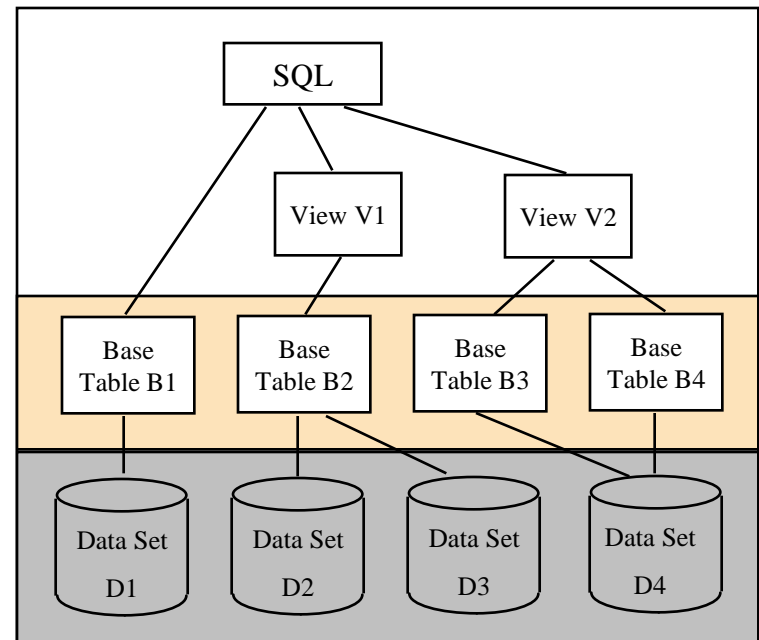


<e.g.2> CREATE INDEX X ON P (P#) CLUSTER;

Note: A given base table can have at most one cluster index.

Views

- Virtual table (doesn't really exist)
- No stored file
- Definition of view is stored in **system catalog**
- A base table may be stored in several files
- A file may contain several base tables
- A view may be derived from several base tables
- A base table may derive several views



Views: Example

<e.g.> **CREATE VIEW LONDON-SUPPLIERS** view name
AS SELECT S# , SNAME , STATUS } view
FROM S } definition
WHERE CITY = ' LONDON' ? } in catalog



LONDON-SUPPLIERS



S#	SNAME	STATUS
S1	Smith	20
S4	Clark	20

S#

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens

Views: Example (cont.)

Can be used as base table: e.g. S, P, SP

<e.g.>

SELECT *

FROM LONDON-SUPPLIERS

WHERE STATUS < 50



converted by Bind (ref. p. 2-7)

SELECT S# , SNAME , STATUS

FROM S

WHERE STATUS < 50

AND CITY = ' LONDON'

Views: Advantages

■ Advantages of views:

(1) Provide logical data independence :

Logical data independence (e.g. Relation): users and user programs are not dependent on logical structure of the database.

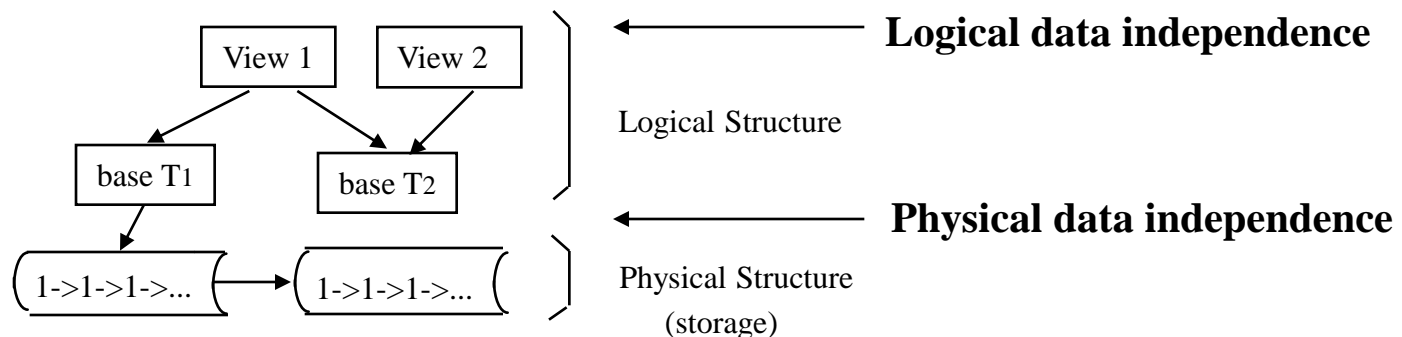
Two aspects of logical data independence : Growth and restructuring.

(v.s. **Physical data independence** (e.g. B-tree): users and user programs are not dependent on physical structure of the stored database.)

(2) Allow same data to be seen by different users in different ways.

(3) User perception is simplified.

(4) Automatic security is provided for hidden data



2.3 Data Manipulation

DML of SQL

- ❑ **Retrieval Operation**
 - SELECT
- ❑ **Update Operation**
 - UPDATE
 - DELETE
 - INSERT
- ❑ **Expressions**
 - Table Expressions
- ❑ **Operations on View**

Retrieval Operations

- Get color and city for "non-Paris" parts with weight greater than ten.

```
SELECT P.COLOR, P.CITY
FROM P
WHERE P.CITY <> 'Paris'
AND P.WEIGHT > 10;
```

P

P#	PNAME	COLOR	WEIGHT	CITY
P1	Nut	Red	12	London
P2	Bolt	Green	17	Paris
P3	Screw	Blue	17	Rome
P4	Screw	Red	14	London
P5	Cam	Blue	12	Paris
P6	Cog	Red	19	London

- DISTINCT

```
SELECT DISTINCT P.COLOR, P.CITY
FROM P
WHERE P.CITY <> 'Paris'
AND P.WEIGHT > 10;
```

P#	COLOR	CITY
P1	Red	London
P3	Blue	Rome
P4	Red	London
P6	Red	London

- ORDER BY

```
SELECT DISTINCT P.COLOR, P.CITY
FROM P
WHERE P.CITY <> 'Paris'
AND P.WEIGHT > 10
ORDER BY CITY DESC;
```

P#	COLOR	CITY
P3	Blue	Rome
P1	Red	London
P4	Red	London
P6	Red	London

Retrieval Operations (cont.)

- For all parts, get the part number and the weight of that part in grams.

```
SELECT P.P#, P.WEIGHT * 454 AS GMWT  
FROM P ;
```

– If the AS MWT specification had been omitted, the corresponding result column would effectively have been unnamed.

- Get full details of all suppliers.

```
SELECT * – or "SELECT S.*" (i.e., the "*" can be qualified )  
FROM S ;
```

- Get the total number of suppliers.

```
SELECT COUNT ( * ) AS N  
FROM S ;
```

Retrieval Operations (cont.)

- Get the maximum and minimum quantity for part P2.

```
SELECT MAX (SP.QTY) AS MAXQ,  
        MIN (SP.QTY) AS MINQ  
FROM    SP  
WHERE   SP.P# = 'P2';
```

- For each part supplied, get the part number and the total shipment quantity.

```
SELECT SP.P#, SUM (SP.QTY) AS TOTQTY  
FROM    SP  
GROUP BY SP.P#;
```


Retrieval Operations (cont.)

- Get part numbers for all parts supplied by more than one supplier.

```
SELECT SP.P#  
FROM SP  
GROUP BY SP.P#  
HAVING COUNT ( SP. S# ) > 1;
```

- Get supplier names for suppliers who supply part P2.

```
SELECT DISTINCT S.SNAME  
FROM S  
WHERE S. S# IN  
      ( SELECT SP. S#  
        FROM SP  
        WHERE SP.P# = 'P2');
```

Update Operations

- Single-row INSERT.

```
INSERT
  INTO  P(P#, PNAME, COLOR, WEIGHT, CITY)
VALUES ('P8', 'Sprocket', 'Pink', 14, 'Nice');
```

- Multi-row INSERT.

```
INSERT
  INTO  TEMP (S#, CITY)
SELECT S.S#, S.CITY
FROM    S
WHERE   S.STATUS > 15 ;
```

- Multi-row UPDATE.

```
UPDATE P
SET    COLOR = 'Yellow',
       WEIGHT = P.WEIGHT + 5
WHERE  P.CITY = 'Paris';
```

Update Operations (cont.)

- Multi-row UPDATE.

```
UPDATE P
SET     CITY = ( SELECT S.CITY
                  FROM S
                  WHERE S.S# = 'S5')
WHERE  P.COLOR = ' Red '
```

- Multi-row DELETE.

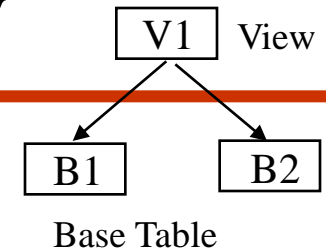
```
DELETE
FROM   SP
WHERE  'London' = (SELECT S.CITY
                  FROM S
                  WHERE S.S# = SP.S#);
```

Table Expressions

- The SELECT, FROM, WHERE, GROUP BY, and HAVING clause.
- A Comprehensive Example

```
SELECT P.P#, 'Weight in grams = ' AS TEXT1,  
       P.WEIGHT * 454 AS GMWT,  
       P.COLOR, 'Max quantity = ' AS TEXT2,  
       MAX (SP.QTY) AS MQY  
FROM   P, SP  
WHERE  P.P# = SP.P#  
AND    ( P.COLOR = 'Red' OR P.COLOR = ' Blue')  
AND    SP.QTY > 200  
GROUP BY P.P#, P.WEIGHT, P.COLOR  
HAVING SUM (SP.QTY) > 350;
```

DML operations on View



- Retrieval (SELECT): no problem
- Update (INSERT, DELETE, UPDATE): ?

(1) Column Subset: theoretically updateable if contains primary key.

<e.g.1> : **CREATE VIEW S#_CITY**
AS SELECT S#, CITY
FROM S;

S

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens

Base Table

S#_CITY

S#	CITY
S1	London
S2	Paris
S3	Paris
S4	London
S5	Athens

View

INSERT
INTO S#_CITY
VALUES ('S6', 'Rome');

S

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens
S6	<i>Null</i>	<i>Null</i>	<i>Rome</i>

DML on View: Column Subset without key

<e.g.2>

```
CREATE VIEW STATUS_CITY
AS SELECT STATUS, CITY
FROM S;
```

S

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens

STATUS_CITY

STATUS	CITY
20	London
10	Paris
30	Paris
30	Athens



```
INSERT
INTO STATUS_CITY
VALUES (30, 'Rome')
```

S

S#	SNAME	STATUS	CITY
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	30	Paris
S4	Clark	20	London
S5	Adams	30	Athens
<u>Null</u>	<u>Null</u>	30	Rome



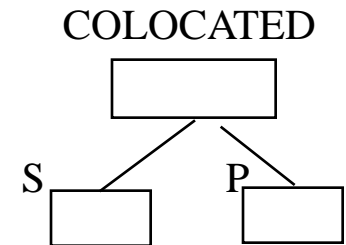
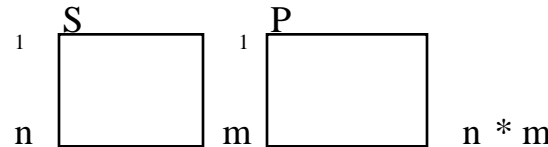
Primary key cannot be null !!

DML on View: Row Subset and Join

(2) Row Subset: *updateable!*

(3) Join: some are not updateable.

- CREATE VIEW COLOCATED (S#, SNAME, S.CITY, P#, PNAME, P.CITY)
AS SELECT S# , SNAME , S.CITY, P# , PNAME , P.CITY
FROM S, P
WHERE S.CITY=P.CITY;



- If issued the following command:

UPDATE COLOCATED

SET S.CITY = 'Athens'

WHERE S.CITY = 'London'

- Then S.CITY ≠ P.CITY



Violate the definition of the view!!

COLOCATED

S#	SNAME	S.CITY	P#	PNAME	P.CITY
S1	Smith	London	P1	Nut	London
S1	Smith	London	P4	Screw	London
S1	Smith	London	P6	Cog	London
S2	Jones	Paris	P2	Bolt	Paris
S2	Jones	Paris	P5	Cam	Paris
S3	Blake	Paris	P2	Bolt	Paris
S3	Blake	Paris	P5	Cam	Paris
S4	Clark	London	P1	Nut	London
S4	Clark	London	P4	Screw	London
S4	Clark	London	P6	Cog	London

DML on View: Statistical Summary

(4) Statistical Summary : not updateable.

<e.g.> : **CREATE VIEW PQ(P# , TOTQTY)**
 AS SELECT P# , SUM(QTY)
 FROM SP
 GROUP BY P# ;



No stored data item for the field "TOTQTY"

SP

S#	P#	QTY
S1	P1	300
S1	P2	200
S1	P3	400
S1	P4	200
S1	P5	100
S1	P6	100
S2	P1	300
S2	P2	400
S3	P2	200
S4	P2	200
S4	P4	300
S4	P5	400

P#	TOTQTY
—	—
P1	600
P2	1000
⋮	⋮

PQ

P#	TOTQTY
P1	600

View

SP

S#	P#	QTY

Base

2.4 The System Catalog

System Catalog: Concept

- **The Catalog Structure**

- SYSTABLES
- SYSCOLUMNS
- SYSINDEXES

	cc	record	creator	...
S	4	5	Yang	
P	5	6	Yang	
SP	3	12	Yang	

- **System Catalog: "Data about data"** [\(Ref. p. 2-8\)](#)

i.e. information of base tables, view, indexes, users, and access privileges that are of interest to the system itself.

- **Optimizer:** use index info. to choose access method. [\(Ref. p. 2-9\)](#)
- **Authorization subsystem:** use access privilege to grant or deny user requests.

	T1	T2	...
u ₁	R	R	
u ₂	W	R	
⋮			
u _n	O	W	

access
control
matrix

- **Querying the catalog:** by SQL DML

Updating the Catalog

- Cannot use the SQL update, delete and insert, because it would be too easy to destroy information!
- It is data definition statements (i.e. CREATE, DROP, ALTER) that perform such updates.
 - CREATE = INSERT into catalog
 - DROP = DELETE from catalog
 - ALTER = UPDATE catalog

Updating the Catalog: Example

- <e.g.>: CREATE TABLE S
 (S# CHAR(5) Not Null,
 SNAME CHAR(20) ,
 STATUS SMALLINT,
 CITY CHAR(5);



SYSTABLE

NAME	CREATOR	
...		
...
...
...
S		

SYSCOLUMNS

NAME	TBNAME	...
...	...	
...	...	
...	...	
S#	S	
SNAME	S	
STATUS	S	
CITY	S	

Updating the Catalog: COMMENT

- Catalog also includes entries for catalog tables.

SYSTABLE

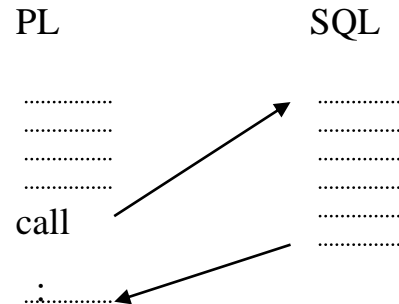
NAME	CREATOR	...	REMARK
SYSTABLE	SYSIBM
SYSCOLUMN	SYSIBM
.
.
S	Yang	...	Supplier
P	Yang	...	Part
SP	Yang

- The only statement that updates catalog: COMMENT
 - <e.g.>: COMMENT ON TABLE S IS Supplier

2.5 Embedded SQL

Embedded SQL: Dual-mode

- Dual-mode principle: any SQL statement that can be used at terminal (interactive), can also be used in an application program (programmable).



- **PL/I** (Record operations) vs. **SQL** (Set operations)

Embedded SQL: a Fragment

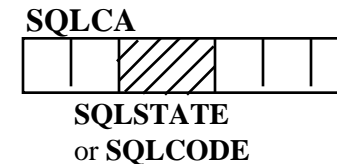
- <e.g.> Fragment of a PL/I program with embedded SQL

```
1 EXEC SQL BEGIN DECLARE SECTION ;
2     DCL SQLSTATE CHAR(5) ;
3     DCL P#          CHAR(6) ;
4     DCL WEIGHT     FIXED DECIMAL(3) ;
5 EXEC SQL END DECLARE SECTION ;
6 P# = ' P2 ' ;          /* for example          */
7 EXEC SQL SELECT P.WEIGHT
8             INTO   :WEIGHT
9             FROM   P
10            WHERE P.P# = :P# ;
11 IF SQLSTATE = ' 00000 '
12 THEN .... ;          /* WEIGHT = retrieved value */
13 ELSE .... ;         /* some exception occurred  */
```


Embedded SQL: a Fragment (cont.)

1. Embedded SQL statements are prefix by EXEC SQL.
2. Executable statements can appear wherever.
(non-executable statements: e.g. DECLARE TABLE, DECLARE CURSOR).
3. SQL statements can reference host variable. (PL/I var., :City)
4. Any table used should be declared by DECLARE TABLE, because it is used by pre-compiler.
5. **SQLSTATE/SQLCODE**: feedback information of SQL, stored in **SQLCA** (SQL Communication Area).

SQLSTATE	= 0	success
	> 0	warning
	< 0	error



6. Host variables must have compatible data type with SQL field.
7. Host variables can have same name as database fields.
e.g. **City**, **:City**
(SQL) (PL/I)

Operation: Singleton SELECT

- Singleton SELECT:

```
EXEC SQL SELECT STATUS  
INTO   :RANK  
FROM   S  
WHERE  S#=: GIVENS#;
```

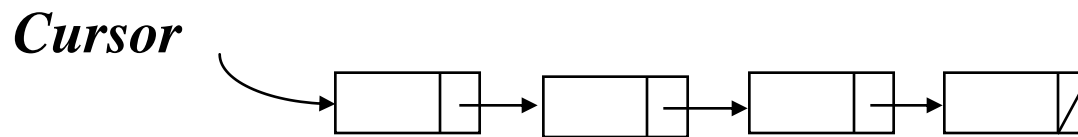
- If only one record is satisfied: `SQLCODE = 0`
- If no record is satisfied: `SQLCODE > 0`
- If more than one record are satisfied: `SQLCODE < 0`
- How to deal with NULL value? Indicator variable!

```
EXEC SQL SELECT STATUS  
INTO   :RANK :RANKIND  
FROM   S  
WHERE  S#=:GIVENS#
```

- RANKIND: an indicator variable, 15-bit signed binary integer.
- If `RANKIND = -1 THEN/* Status was NULL */`

Operation: Multiple SELECT

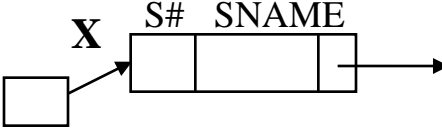
- Multiple SELECT:
 - How to handle the cases that more than one record are satisfied?



Cursor

- A kind of pointer that can be run through a set of records.

define cursor { EXEC SQL DECLARE **X** CURSOR FOR /*define cursor S*/
 SELECT S#, SNAME
 FROM S
 WHERE CITY =:Y;

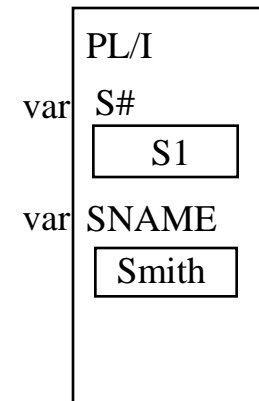
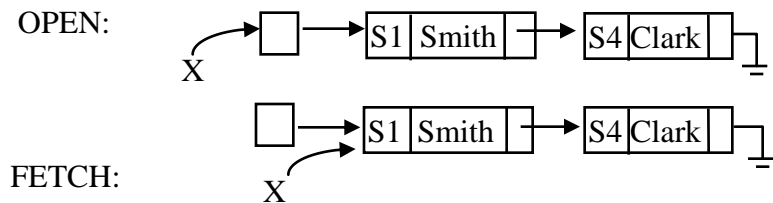


The diagram shows a small square box labeled 'X' with an arrow pointing to a larger rectangular box representing a record. This record box is divided into two sections: the left section is labeled 'S#' and the right section is labeled 'SNAME'. An arrow points from the right side of the 'SNAME' section to the right, indicating the next record in the sequence.

open cursor { EXEC SQL OPEN **X**; /*activate cursor, execute the query*/
 DO for all S records accessible via X;
 EXEC SQL FETCH **X** INTO :S#, :SNAME
 /*advance pt., assign values */
 END;

close cursor { EXEC SQL CLOSE **X**; /*deactivate cursor X*/

e.g. Y = 'London'



Embedded SQL: An Example

- Embedded SQL

A comprehensive example

The program accepts four input values : a part number (GIVENP#), a city name (GIVENCIT), a status increment (GIVENINC), and a status level (GIVENLVL). The program scans all suppliers of the part identified by GIVENP#. For each such supplier, if the supplier city is GIVENCIT, then the status is increased by GIVENINC; otherwise, if the status is less than GIVENLVL, the supplier is deleted, together with all shipments for that supplier. In all cases supplier information is listed on the printer, with an indication of how that particular supplier was handled by the program.

Embedded SQL: An Example (cont.)

SQLEX: PROC OPTIONS (MAIN) ;

```
DCL GIVENP #          CHAR(6) ;  
DCL GIVENCIT         CHAR(15) ;  
DCL GIVENINC         FIXED BINARY(15) ;  
DCL GIVENLVL        FIXED BINARY(15) ;  
DCL S#               CHAR(5) ;  
DCL SNAME            CHAR(20) ;  
DCL STATUS           FIXED BINARY(15) ;  
DCL CITY             CHAR(15) ;  
DCL DISP             CHAR(7) ;  
DCL MORE_SUPPLIERS  BIT(1) ;
```

} PL/I Var.

```
EXEC SQL INCLUDE SQLCA ; /* p.2-41 */
```

```
EXEC SQL DECLARE S TABLE
```

```
( S#          CHAR(5)  NOT NULL,  
  SNAME      CHAR(20) NOT NULL,  
  STATUS     SMALLINT NOT NULL,  
  CITY       CHAR(20) NOT NULL ) ;
```

```
EXEC SQL DECLARE SP TABLE
```

```
( S#          CHAR(5)  NOT NULL,  
  P#          CHAR(6)  NOT NULL,  
  QTY         INTEGER  NOT NULL ) ;
```

[Back 2-7](#)

Embedded SQL: An Example (cont.)

```
EXEC SQL DECLARE Z CURSOR FOR
      SELECT S#, SNAME, STATUS, CITY
      FROM   S
      WHERE  EXISTS
            ( SELECT *
              FROM   SP
              WHERE  SP.S# = S.S#
              AND    SP.P# = :GIVENP# )   i.e. P2
      FOR UPDATE OF STATUS ;
```

```
EXEC SQL WHENEVER NOT FOUND CONTINUE ;
EXEC SQL WHENEVER SQLERROR CONTINUE ;
EXEC SQL WHENEVER SQLWARNING CONTINUE ;
```

```
ON CONDITION ( DBEXCEPTION )
```

```
BEGIN ;
```

```
    PUT SKIP LIST ( SQLCA ) ;
```

```
    EXEC SQL ROLLBACK ;
```

```
    GO TO QUIT ;
```

```
END ;
```

Embedded SQL: An Example (cont.)

```
Main → GET LIST ( GIVENP#, GIVENCIT, GIVENINC, GIVENLVL ) ;
        EXEC SQL OPEN Z ;
        IF SQLCODE <> 0  /* abnormal */
        THEN SIGNAL CONDITION ( DBEXCEPTION ) ;
SQLCODE =0 → MORE_SUPPLIERS = ' 1' B ;
        DO WHILE ( MORE_SUPPLIERS ) ;
            EXEC SQL FETCH Z INTO :S#, :SNAME, :STATUS, :CITY ;
            SELECT ; /* case */ /* a PL/I SELECT, not a SQL SELECT */
            WHEN ( SQLCODE = 100 ) /* Not found */
                MORE_SUPPLIERS = ' 0' B ;
            WHEN ( SQLCODE <> 100 & SQLCODE <> 0 ) /* Warning */
                SIGNAL CONDITION ( DBEXCEPTION ) ;
```


Embedded SQL: An Example (cont.)

```
WHEN ( SQLCODE = 0 ) /* success */
    DO ;
    DISP = ' bbbbbbb ' ; /* empty the display buffer */
    IF CITY = GIVENCIT
    THEN
        DO ;
            EXEC SQL UPDATE S
                SET STATUS = STATUS + : GIVENINC ;
                WHERE CURRENT OF Z ;
            IF SQLCODE <> 0
            THEN SIGNAL CONDITION ( DBEXCEPTION ) ;
            DISP = ' UPDATED ' ;
        END ;
    ELSE
        IF STATUS < GIVENLVL
        THEN
            DO ;
                EXEC SQL DELETE
                    FROM SP
                    WHERE S# = : S# ;
```

Embedded SQL: An Example (cont.)

```
IF SQLCODE <> 0 & SQLCODE <> 100
THEN SIGNAL CONDITION ( DBEXCEPTION );
EXEC SQL DELETE
        FROM S
        WHERE CURRENT OF Z ;
IF SQLCODE <> 0
THEN SIGNAL CONDITION ( DBEXCEPTION);
DISP = 'DELETED' ;
END ;
        PUT SKIP LIST ( S#, SNAME, STATUS, CITY, DISP ) ;
END ; /* WHEN ( SQLCODE = 0 ) */
END ; /* PL/I SELECT */
END ; /* DO WHILE */
EXEC SQL CLOSE Z ;
EXEC SQL COMMIT ; /* normal termination */
QUIT: RETURN ;
END ; /* SQLEX */
```

Program Exercise 1: Using DBMS

- ❑ Due Date:
- ❑ EX. 4.1-4.6 (p.99-100)

Using the suppliers-parts-projects database, write a program with embedded SQL statements to list all supplier rows, in supplier number order. Each supplier row should be immediately followed in the listing by all project rows for projects supplied by that supplier, in project number order.

- create database
- selection
- update
- query catalog
- .
- .
- .
- embedded SQL (program)